

A Proof of Lemma 3.6

In this section, we present a comprehensive proof of Lemma 3.6. This lemma establishes a fundamental property of our recursive construction, as defined in Definition 3.2: the only ϵ^2 -approximate Nash equilibrium is located at the maximum value element. Consequently, the fictitious play dynamics fail to reach even an approximate Nash equilibrium (within the specified degree of approximation) unless the strategy $(\frac{n}{2}, \frac{n}{2} + 1)$ has been played for a sufficient duration. Our theoretical findings are further supported by our experiments. As shown in Figure 4b, the Nash gap does not vanish before the final strategy switch occurs. To provide a precise definition of the Nash gap, we present it below.

Definition A.1 (Nash Gap for Identical Payoff). The Nash gap at round t for a two-player identical payoff games is

$$\left(\max_{i \in [n]} [Ay^{(t)}]_i - (x^{(t)})^\top Ay^{(t)} \right) + \left(\max_{j \in [n]} [(x^{(t)})^\top A]_j - (x^{(t)})^\top Ay^{(t)} \right)$$

To establish Lemma 3.6, we employ a similar approach as in the other proofs presented in this work. Specifically, we heavily rely on the structure of our payoff matrix and utilize an induction technique to demonstrate that the majority of the probability mass must be concentrated in the maximum element. The induction argument starts from the outermost elements of the matrix, namely row $n-1$, column 1, and row 1, column $n-1$. By successive induction steps, we establish that until the maximum element is reached, none of the elements in those rows or columns can possess a significant probability mass.

Lemma A.2 (Proof of Lemma 3.6). *Let $\epsilon \in (0, 1/56n^3]$ and (x^*, y^*) an ϵ^2 -NE. Then for each $i \in [0, n/2 - 2]$,*

$$\bullet \quad x_{i+1}^* \leq \epsilon \text{ and } x_{n-i}^* \leq \epsilon.$$

$$\bullet \quad y_{i+1}^* \leq \epsilon \text{ and } y_{n-i}^* \leq \epsilon.$$

Proof. We will prove the claim by induction. First, assume that for all $j \in [0, i-1]$, we have:

$$\bullet \quad x_{j+1}^* \leq \epsilon \text{ and } x_{n-j}^* \leq \epsilon.$$

$$\bullet \quad y_{j+1}^* \leq \epsilon \text{ and } y_{n-j}^* \leq \epsilon.$$

Next, we proceed to establish the inequalities $x_{n-i}^*, y_{i+1}^*, x_{i+1}^*, y_{n-i}^* \leq \epsilon$. We demonstrate these inequalities in the exact order as presented, as their proof relies on the underlying structure of the matrix. It is important to note that there are interdependencies between these inequalities, which we will address accordingly.

Case 1: $x_{n-i}^* \leq \epsilon$

Let assume that $x_{n-i}^* > \epsilon$ and we will reach a contradiction. From Observation 3.3, we notice that the utility of row $n-i$ equals

$$[Ay^*]_{n-i} = (4i+1)y_{i+1}^* + 4iy_{n-i+1}^* \tag{4}$$

At the same time the utility of row $i+1$ equals

$$[Ay^*]_{i+1} = (4i+2)y_{i+1}^* + (4i+3)y_{n-i}^*$$

As a result, by taking the difference on the utilities of row $i+1$ and $n-i$ we get,

$$\begin{aligned} [Ay^*]_{i+1} - [Ay^*]_{n-i} &= (4i+2)y_{i+1}^* + (4i+3)y_{n-i}^* - (4i+1)y_{i+1}^* - 4iy_{n-i+1}^* \\ &= y_{i+1}^* + (4i+3)y_{n-i}^* - 4iy_{n-i+1}^* \\ &\geq y_{i+1}^* + y_{n-i}^* - 4i\epsilon \end{aligned}$$

where the last inequality follows by the fact that $y_{n-i+1}^* \leq \epsilon$ (Inductive Hypothesis). As a result, we conclude that

$$[Ay^*]_{i+1} - [Ay^*]_{n-i} \geq y_{i+1}^* + y_{n-i}^* - 4i\epsilon$$

434 In case $y_{i+1}^* + y_{n-i}^* \geq (4i+1)\epsilon$ then $[Ay^*]_{i+1} - [Ay^*]_{n-i} \geq \epsilon$. Hence if the row player puts x_{n-i}^*
 435 probability mass to row $i+1$ by transferring the probability mass from row $n-i$ to row $i+1$ then it
 436 increases its payoff by $x_{n-i}^*([Ay^*]_{i+1} - [Ay^*]_{n-i}) > \epsilon^2$. The latter contradicts with the assumption
 437 that (x^*, y^*) is an ϵ^2 -NE. Thus, we conclude in the following two statements,

$$y_{i+1}^* + y_{n-i}^* \leq (4i+1)\epsilon \quad \text{and} \quad [Ay^*]_{n-i} \leq 2(4i+1)^2\epsilon$$

438 where the last inequality is obtained by combining the first inequality with Equation (4). Now
 439 consider the sum of the utilities of rows $k \in [i+2, n-i-1]$. By the construction of the payoff
 440 matrix A we can easily establish the following claim.

441 **Proposition A.3.** *The sum of utilities of rows $k \in [i+2, n-i-1]$ satisfies the inequality,*

$$\sum_{k=i+2}^{n-i-1} [Ay^*]_k \geq \sum_{k=i+2}^{n-i-1} y_k^*$$

442 By Proposition A.3 we are ensured that

$$\begin{aligned} \sum_{k=i+2}^{n-i-1} [Ay^*]_k &\geq \sum_{k=i+2}^{n-i-1} y_k^* \\ &= 1 - \sum_{k=1}^{i+1} y_k^* - \sum_{k=n-i}^n y_k^* \\ &= 1 - (y_{i+1}^* + y_{n-i}^*) - \left(\sum_{k=1}^i y_k^* + \sum_{k=n-i+1}^n y_k^* \right) \\ &\geq 1 - (4i+1)\epsilon - n\epsilon \\ &\geq 1 - (5n+1)\epsilon \end{aligned}$$

443 where the second inequality follows by the fact that $y_{i+1}^* + y_{n-i}^* \leq (4i+1)\epsilon$ and the fact that $y_j^* \leq \epsilon$
 444 for all $k \in [1, i] \cup [n-i+1, n]$ (Inductive Hypothesis).

445 Due to the fact that

$$\sum_{k=i+2}^{n-i-1} [Ay^*]_k \geq 1 - (5n+1)\epsilon$$

446 we are ensured that there exists a row $k^* \in [i+2, n-i-1]$ with utility $[Ay^*]_{k^*} \geq \frac{1-(5n+1)\epsilon}{n}$. Now
 447 consider the difference between the utility of row k^* and the row $n-i$.

$$[Ay^*]_{k^*} - [Ay^*]_{n-i} \geq \frac{1-(5n+1)\epsilon}{n} - 2(4i+1)^2\epsilon \geq \frac{1-(5n+1)\epsilon}{n} - 2(4n+1)^2\epsilon \geq \epsilon$$

448 where the last inequality holds for $\epsilon \leq 1/56n^3$. Hence if the row player puts x_{n-i}^* probability mass
 449 to row k^* then it increases its payoff by $x_{n-i}^*([Ay^*]_{k^*} - [Ay^*]_{n-i}) > \epsilon^2$. The latter contradicts
 450 with the assumption that (x^*, y^*) is an ϵ^2 -NE. Thus we have reached to a final contradiction that
 451 $x_{n-i}^* > \epsilon$.

452 **Case 2:** $y_{i+1}^* \leq \epsilon$

453 Similar to the previous case, we assume that $y_{i+1}^* > \epsilon$ and proceed to derive a contradiction. From
 454 Observation 3.3, we notice that the utility of column $i+1$ is given by:

$$[(x^*)^\top A]_{i+1} = (4i+2)x_{i+1}^* + (4i+1)x_{n-i}^*$$

455 At the same time the utility of column $n-i$ equals

$$[(x^*)^\top A]_{n-i} = (4i+3)x_{i+1}^* + (4i+4)x_{n-i-1}^*$$

456 As a result, by taking the difference on the utilities of columns $n-i$ and $i+1$ we get,

$$\begin{aligned}
[(x^*)^\top A]_{n-i} - [(x^*)^\top A]_{i+1} &= (4i+3)x_{i+1}^* + (4i+4)x_{n-i-1}^* - (4i+2)x_{i+1}^* - (4i+1)x_{n-i}^* \\
&= x_{i+1}^* + (4i+4)x_{n-i-1}^* - (4i+1)x_{n-i}^* \\
&\geq x_{i+1}^* + x_{n-i-1}^* - (4i+1)\epsilon
\end{aligned}$$

457 where the last inequality follows by the fact $x_{n-i}^* \leq \epsilon$ (Inductive Step Case 1). As a result, we
458 conclude that

$$[(x^*)^\top A]_{n-i} - [(x^*)^\top A]_{i+1} \geq x_{i+1}^* + x_{n-i-1}^* - (4i+1)\epsilon$$

459 In case $x_{i+1}^* + x_{n-i-1}^* \geq (4i+2)\epsilon$ then $[(x^*)^\top A]_{n-i} - [(x^*)^\top A]_{i+1} \geq \epsilon$. Hence if the column
460 player puts y_{i+1}^* probability mass to column $n-i$ then it increases its payoff by $y_{i+1}^*([(x^*)^\top A]_{n-i} -$
461 $[(x^*)^\top A]_{i+1}) > \epsilon^2$. The latter contradicts with the assumption that (x^*, y^*) is an ϵ^2 -NE. Thus we
462 conclude in the following two statements:

$$x_{i+1}^* + x_{n-i-1}^* \leq (4i+2)\epsilon \quad \text{and} \quad [(x^*)^\top A]_{i+1} \leq 2(4i+2)^2\epsilon$$

463 Now consider the sum of the utilities of columns $k \in [i+2, n-i-1]$. By the construction of the
464 payoff matrix A we can easily establish the following claim.

465 **Proposition A.4.** *The sum of utilities of columns $k \in [i+2, n-i-1]$ satisfies the inequality,*

$$\sum_{k=i+2}^{n-i-1} [(x^*)^\top A]_k \geq \sum_{k=i+2}^{n-i-1} x_k^*$$

466 By Proposition A.4 we are ensured that

$$\begin{aligned}
\sum_{k=i+2}^{n-i-1} [(x^*)^\top A]_k &\geq \sum_{k=i+2}^{n-i-1} x_k^* \\
&= 1 - \sum_{k=1}^{i+1} x_j^* - \sum_{k=n-i}^n x_j^* \\
&= 1 - (x_{i+1}^* + x_{n-i}^*) - \left(\sum_{k=1}^i x_i^* + \sum_{k=n-i+1}^n x_j^* \right) \\
&\geq 1 - ((4i+2)\epsilon + \epsilon) - \left(\sum_{k=1}^i x_i^* + \sum_{k=n-i+1}^n x_j^* \right) \\
&\geq 1 - (5n+3)\epsilon
\end{aligned}$$

467 where the second to last inequality follows by the facts: $x_{i+1}^* + x_{n-i-1}^* \leq (4i+2)\epsilon$ and so
468 $x_{i+1}^* \leq (4i+2)\epsilon$, and the Inductive Step Case 1, $x_{n-i}^* \leq \epsilon$. Moreover, the last inequality holds by
469 the Inductive Hypothesis: $x_k^* \leq \epsilon$ for all $k \in [1, i] \cup [n-i+1, n]$.

470 Due to the fact that

$$\sum_{k=i+2}^{n-i-1} [(x^*)^\top A]_k \geq 1 - (5n+3)\epsilon$$

471 we are ensured that there exists a column $k^* \in [i+2, n-i-1]$ with utility $[(x^*)^\top A]_{k^*} \geq \frac{1-(5n+3)\epsilon}{n}$.
472 Now consider the difference between the utility of column k^* and the column $i+1$.

$$[(x^*)^\top A]_{k^*} - [(x^*)^\top A]_{i+1} \geq \frac{1-(5n+3)\epsilon}{n} - 2(4i+2)^2\epsilon \geq \frac{1-(5n+3)\epsilon}{n} - 2(4n+2)^2\epsilon \geq \epsilon$$

where the last inequality follows by the fact that $\epsilon \leq 1/56n^3$. Hence if the column player puts y_{i+1}^* probability mass to column k^* then it increases its payoff by $y_{i+1}^*([(x^*)^\top A]_{k^*} - [(x^*)^\top A]_{i+1}) > \epsilon^2$. The latter contradicts with the assumption that (x^*, y^*) is an ϵ^2 -NE. Thus we have reached to a final contradiction that $y_{i+1}^* > \epsilon$.

Case 3: $x_{i+1}^* \leq \epsilon$

Let assume that $x_{i+1}^* > \epsilon$ and we will reach a contradiction. From Observation 3.3, we notice that the utility of row $i + 1$ equals,

$$[Ay^*]_{i+1} = (4i + 2)y_{i+1}^* + (4i + 3)y_{n-i}^*$$

Next, we examine the row $n - (i + 1)$ in the inner submatrix. If this row is not well-defined, it implies that the inductive step $j = i$ has reached the 2×2 submatrix. In this case, the proof of Lemma 3.6 has already been completed. Otherwise, the utility of row $n - (i + 1)$ equals

$$[Ay^*]_{n-(i+1)} = (4(i + 1) + 1)y_{(i+1)+1}^* + (4(i + 1))y_{n-(i+1)+1}^* \geq (4i + 4)y_{n-i}^*$$

As a result, by taking the difference on the utilities of row $i + 1$ and $n - (i + 1)$ we get,

$$\begin{aligned} [Ay^*]_{n-(i+1)} - [Ay^*]_{i+1} &= (4i + 4)y_{n-i}^* - (4i + 2)y_{i+1}^* - (4i + 3)y_{n-i}^* \\ &= y_{n-i}^* - (4i + 2)y_{i+1}^* \\ &\geq y_{n-i}^* - (4i + 2)\epsilon \end{aligned}$$

where the last inequality follows by the fact that $y_{i+1}^* \leq \epsilon$ (Inductive Step Case 2). As a result, we conclude that

$$[Ay^*]_{n-(i+1)} - [Ay^*]_{i+1} \geq y_{n-i}^* - (4i + 2)\epsilon$$

In case $y_{n-i}^* \geq (4i + 3)\epsilon$ then $[Ay^*]_{n-(i+1)} - [Ay^*]_{i+1} \geq \epsilon$. Hence, if the row player puts x_{i+1}^* probability mass to row $n - (i + 1)$ then it increases its payoff by $x_{i+1}^*([Ay^*]_{n-(i+1)} - [Ay^*]_{i+1}) > \epsilon^2$. The latter contradicts with the assumption that (x^*, y^*) is an ϵ^2 -NE. Thus, we conclude the following two statements:

$$y_{n-i}^* \leq (4i + 3)\epsilon \quad \text{and} \quad [Ay^*]_{i+1} \leq 2(4i + 3)^2\epsilon$$

Now consider the sum of the utilities of rows $k \in [i + 2, n - i - 1]$. From Proposition A.3

$$\begin{aligned} \sum_{k=i+2}^{n-i-1} [Ay^*]_k &= \sum_{k=i+2}^{n-i-1} y_k^* = 1 - \left(\sum_{k=1}^{i+1} y_k^* + \sum_{k=n-i}^n y_k^* \right) \geq 1 - y_{n-i}^* - \left(\sum_{k=1}^{i+1} y_k^* + \sum_{k=n-i+1}^n y_k^* \right) \\ &\geq 1 - (4i + 3)\epsilon - n\epsilon = 1 - (5n + 3)\epsilon \end{aligned}$$

Thus, we are ensured that there exists a row $k^* \in [i + 2, n - i - 1]$ with utility $[Ay^*]_{k^*} \geq \frac{1 - (5n + 3)\epsilon}{n}$. Now consider the difference between the utility of row k^* and the row $n - i$.

$$[Ay^*]_{k^*} - [Ay^*]_{i+1} \geq \frac{1 - (5n + 3)\epsilon}{n} - 2(4i + 3)^2\epsilon \geq \frac{1 - (5n + 3)\epsilon}{n} - 2(4n + 3)^2\epsilon \geq \epsilon$$

where the last inequality follows by the fact that $\epsilon \leq 1/56n^3$. Hence if the row player puts x_{i+1}^* probability mass to row k^* then it increases its payoff by $x_{i+1}^*([Ay^*]_{k^*} - [Ay^*]_{i+1}) > \epsilon^2$. The latter contradicts with the assumption that (x^*, y^*) is an ϵ^2 -NE. Thus we have reached to a final contradiction that $x_{i+1}^* > \epsilon$.

Case 4: $y_{n-i}^* \leq \epsilon$

Let assume that $y_{n-i}^* > \epsilon$ and we will reach a contradiction. From Observation 3.3, we notice that the utility of column $n - i$ equals,

$$[(x^*)^\top A]_{n-i} = (4i + 3)x_{i+1}^* + (4i + 4)x_{n-(i+1)}^*$$

Now, we consider the column $(i + 1) + 1$ in the inner submatrix. In case this column is not well-defined, it means that the inductive step $j = i$ has reached the 2×2 submatrix and so the proof of the Lemma 3.6 has already been completed. Otherwise, the utility of column $(i + 1) + 1$ equals

$$[(x^*)^\top A]_{(i+1)+1} = (4(i + 1) + 2)x_{(i+1)+1}^* + (4(i + 1) + 1)x_{n-(i+1)+1}^* \geq (4i + 5)x_{n-i}^*$$

As a result, by taking the difference on the utilities of columns $i + 1$ and $n - (i + 1)$ we get,

$$\begin{aligned} [(x^*)^\top A]_{(i+1)+1} - [(x^*)^\top A]_{n-i} &= (4i + 5)x_{n-i}^* - (4i + 3)x_{i+1}^* - (4i + 4)x_{n-(i+1)}^* \\ &= x_{n-i}^* - (4i + 3)x_{i+1}^* \\ &\geq x_{n-i}^* - (4i + 3)\epsilon \end{aligned}$$

where the last inequality follows by the fact that $x_{i+1}^* \leq \epsilon$ (Inductive Step Case 3). As a result, we conclude that

$$[(x^*)^\top A]_{(i+1)+1} - [(x^*)^\top A]_{n-i} \geq x_{n-i}^* - (4i + 3)\epsilon$$

In case $x_{n-i}^* \geq (4i + 4)\epsilon$ then $[(x^*)^\top A]_{(i+1)+1} - [(x^*)^\top A]_{n-i} \geq \epsilon$. Hence if the column player puts y_{n-i}^* probability mass to column $(i + 1) + 1$ then it increases its payoff by $y_{n-i}^*([(x^*)^\top A]_{(i+1)+1} - [(x^*)^\top A]_{n-i}) > \epsilon^2$. The latter contradicts with the assumption that (x^*, y^*) is an ϵ^2 -NE. Thus we conclude in the following two statements:

$$x_{n-i}^* \leq (4i + 4)\epsilon \quad \text{and} \quad [(x^*)^\top A]_{n-i} \leq 2(4i + 4)^2\epsilon$$

Now consider the sum of the utilities of columns $k \in [i + 2, n - i - 1]$. From Proposition A.4

$$\begin{aligned} \sum_{k=i+2}^{n-i-1} [(x^*)^\top A]_k &= \sum_{k=i+2}^{n-i-1} x_k^* = 1 - \left(\sum_{k=1}^{i+1} x_k^* + \sum_{k=n-i}^n x_k^* \right) \\ &\geq 1 - x_{n-i}^* - \left(\sum_{k=1}^{i+1} x_k^* + \sum_{k=n-i+1}^n x_k^* \right) \\ &\geq 1 - (4i + 4)\epsilon - n\epsilon = 1 - (5n + 4)\epsilon \end{aligned}$$

Thus, we are ensured that there exists a column $k^* \in [i + 2, n - i - 1]$ with utility $[(x^*)^\top A]_{k^*} \geq \frac{1 - (5n + 4)\epsilon}{n}$. Now consider the difference between the utility of column k^* and the column $n - i$.

$$[(x^*)^\top A]_{k^*} - [(x^*)^\top A]_{n-i} \geq \frac{1 - (5n + 4)\epsilon}{n} - 2(4i + 4)^2\epsilon \geq \frac{1 - (5n + 4)\epsilon}{n} - 2(4n + 4)^2\epsilon \geq \epsilon$$

where the last inequality follows by the fact that $\epsilon \leq 1/56n^3$. Hence if the column player puts y_{n-i}^* probability mass to column k^* then it increases its payoff by $y_{n-i}^*([(x^*)^\top A]_{k^*} - [(x^*)^\top A]_{n-i}) > \epsilon^2$. The latter contradicts with the assumption that (x^*, y^*) is an ϵ^2 -NE. Thus we have reached to a final contradiction that $y_{n-i}^* > \epsilon$.

517

□

518 A.1 Proof of Theorem A.3

519 **Proposition A.5.** *The sum of utilities of rows $k \in [i + 2, n - i - 1]$ admits,*

$$\sum_{k=i+2}^{n-i-1} [Ay^*]_k \geq \sum_{k=i+2}^{n-i-1} y_k^*$$

520 *Proof.* By Observation 3.3 we derive the following equation.

$$[Ay^*]_{i+2} = [Ay^*]_{(i+1)+1} = (4i+2)y_{(i+1)+1}^* + (4i+3)y_{n-(i+1)}^* \geq y_{(i+1)+1}^*$$

521 The claim can be immediately derived from the inequality given above, $\sum_{k=i+2}^{n-i-1} [Ay^*]_k \geq \sum_{k=i+2}^{n-i-1} y_k^*$.

522 □

523 A.2 Proof of Theorem A.4

524 **Proposition A.6.** *The sum of utilities of columns $k \in [i + 2, n - i - 1]$ admits,*

$$\sum_{k=i+2}^{n-i-1} [(x^*)^\top A]_k \geq \sum_{k=i+2}^{n-i-1} x_k^*$$

525 *Proof.* By Observation 3.3 we derive the following equation.

$$[(x^*)^\top A]_{i+2} = [(x^*)^\top A]_{(i+1)+1} \geq (4i+2)x_{(i+1)+1}^* + (4i+3)x_{n-(i+1)}^* \geq x_{(i+1)+1}^*$$

526 The claim can be immediately derived from the inequality given above, $\sum_{k=i+2}^{n-i-1} [(x^*)^\top A]_k \geq \sum_{k=i+2}^{n-i-1} x_k^*$.

527 □

528 B Proof of Lemma 3.8

529 B.1 Omitted Proofs of Section 3.3

530 **Proposition B.1** (Proof of Proposition 3.10). *Let $(i^{(t)}, j^{(t)})$ be a strategy selected by fictitious play*
531 *at round t , and $(i^{(t)}, j^{(t)}) \neq (\frac{n}{2}, \frac{n}{2})$. Then, in a subsequent round, fictitious play will choose the*
532 *strategy of greater value that is either on row $i^{(t)}$ or column $j^{(t)}$.*

533 *Proof.* To establish the claim, we employ the concept of a cumulative utility vector, as defined in
534 Definition 2.7. According to Proposition 3.5, the row $i^{(t)}$ and column $j^{(t)}$ combined have three
535 distinct non-zero elements. Without loss of generality, let's assume that the greater element is in
536 column $j^{(t)}$, but in a different row, denoted as i' .

537 Firstly, we observe that any subsequent strategy will involve only those three elements. This is
538 because in the cumulative utility vector, which determines the strategy to be played in each round,
539 only the coordinates corresponding to those elements are updated as long as the strategy $(i^{(t)}, j^{(t)})$ is
540 being played.

541 Moreover, we can demonstrate that among these elements, the one with the greater value will be
542 played next, and this transition is deterministic. This means that the row player will choose the
543 strategy associated with the greater element. We note that this decision is implicitly affected by the
544 strategy of column player.

Let's first exclude the case where the next strategy switch involves the column player. We initially assumed that the greatest element is in column $j^{(t)}$ but on a different row. Consequently, the only non-zero element in a different column than $j^{(t)}$ must have a smaller utility compared to the element $(i^{(t)}, j^{(t)})$. Therefore, there is no incentive to switch to a strategy with lower utility. This confirms that the column player will not opt for a different strategy, ensuring that the next strategy switch, if it occurs, will necessarily involve the row player.

Regarding the row player, we aim to prove that there will be a round where they will change to a different strategy, and consequently, to a different row. To analyze this, let's examine how the cumulative utility vector of the row player changes from round to round.

$$i := i^{(t)} \in \operatorname{argmax}_{i \in [n]} \left[\sum_{s=1}^{t-1} Ae_{j(s)} \right]_i$$

Therefore, as long as the column player continues to use the same strategy, the row $Ae_{j^{(t)}}$ will repeatedly be added to the cumulative vector, reinforcing the coordinate of row i' . However, this cannot happen indefinitely, as the cumulative utility vector is bounded. After a certain number of rounds, the row player will eventually choose the strategy associated with row i' . This proves that claim for the case of the row player.

Similarly, if the greater element is located in the same row but on a different column, a similar argument can be made to prove that the column player will switch strategies next.

□

Corollary B.2 (Proof of Corollary 3.11). *Let t be a round in which a player changes their strategy. Then exactly one of the following statements is true:*

1. *If the row player changes their strategy at round t , i.e. $i^{(t)} \neq i^{(t-1)}$, then the column player can only make the next strategy switch.*
2. *If the column player changes their strategy at round t , i.e. $j^{(t)} \neq j^{(t-1)}$, then the row player can only make the next strategy switch.*

Proof. This corollary is a simple application of Proposition 3.10. We will only prove the first claim. Let t be the round when the row player changes their strategy, i.e., $i^{(t)} \neq i^{(t-1)}$. According to Proposition 3.5, there are three non-zero elements combined in $i^{(t)}, j^{(t)}$. Since the row player changes their strategy, it follows from Proposition 3.10 that the other element in column $j^{(t)}$ but not in row $i^{(t)}$ must necessarily have a smaller value than $(i^{(t)}, j^{(t)})$. Therefore, the element with the greater value must necessarily be in a different column. Hence, by applying Proposition 3.10, we conclude that the column player can only make the next strategy switch. This proves the claim. □

B.2 Auxiliary Propositions for Lemma 3.6

In this subsection, we provide the full version of the proposition used to establish the proof Lemma 3.6.

Proposition B.3. *There exists a round $T_i^1 > T_i^0$ at which*

- (i) *the strategy profile is $(i+1, i+1)$ for the first time,*
- (ii) *for all rounds $t \in [T_i^0, T_i^1 - 1]$, the strategy profile is $(n-i, i+1)$,*
- (iii) *column $i+1$ admits cumulative utility $C_{i+1}^{(T_i^1)} \geq (4i+1) \cdot (R_{i+1}^{(T_i^0)} + 1)$,*
- (iv) *all rows $k \in [(i+1)+1, n-i-1]$ admit $R_k^{(T_i^1)} = 0$ and all columns $k \in [i+2, n-i]$ admit $C_k^{(T_i^1)} = 0$.*

Proof. The proposition is composed of multiple parts, each of which is proven separately. To begin with, we must establish that the new strategy profile chosen by fictitious play will be $(i+1, i+1)$.

585 According to the inductive hypothesis in Theorem 3.1, we know that the strategy $(n - i, i + 1)$ was
 586 played at round T_i^0 for the first time, while any strategy involving a row $\in [i + 1, n - (i + 1)]$ has not
 587 been played until that point. Additionally, the inductive hypothesis also states that the strategy played
 588 before time T_i^0 was $(n - i, n - i + 1)$. Therefore, according to Corollary 3.11, it is guaranteed that
 589 the next strategy switch will be initiated by the row player.

590 In other words, as long as the column player continues to play their current strategy, strategy
 591 $(n - i, i + 1)$ will be played in every subsequent round, which establishes Item (ii). This implies that
 592 the row player's cumulative utility will increase by Ae_{i+1} . As per Observation 3.3, column $i + 1$
 593 only has non-zero elements in positions $i + 1$ and $n - i$.

$$Ae_{i+1} = [0, \dots, \underbrace{4i + 2}_{i+1}, 0, \dots, 0, \underbrace{4i + 1}_{n-i}, 0, \dots] \quad (5)$$

594 Since strategy $i + 1$ has a higher value, a strategy switch in a later time step is certain. Consequently,
 595 there will be a round T_i^1 in which the strategy $(i + 1, i + 1)$ will be played for the first time, establishing
 596 Item (i).

597 We must determine the point at which the strategy switch to $(i + 1, i + 1)$ will take place. According
 598 to Equation (5), the value of strategy $i + 1$ is exactly one greater than the value of strategy $n - i$.
 599 From the inductive hypothesis, we know that row $i + 1$ has a cumulative utility of zero, whereas
 600 row $n - i$ has a cumulative utility of $R_{i+1}^{(T_i^0)}$. Therefore, it will take precisely $R_{i+1}^{(T_i^0)}$ steps for those
 601 strategies to have equal cumulative utility values. Consequently, the strategy switch will either occur
 602 in that round or the immediate next, as the cumulative utility of row $i + 1$ will have surpassed that of
 603 row $n - i$.

604 In order to proceed with Item (iii), we compute the updated cumulative utility of column $(i + 1)$. As
 605 shown in Equation (5), column $i + 1$ has a value of $(4i + 1)$ at position $n - i$. Thus, if row $n - i$ has
 606 been played for a minimum of $R_{i+1}^{(T_i^0)}$ rounds, then the cumulative utility of column $i + 1$ is greater
 607 than $(4i + 1) \cdot R_{i+1}^{(T_i^0)}$. Given that row $n - i$ has also been played previously (inductive hypothesis), it
 608 is reasonable to conclude that:

$$C_{i+1}^{(T_i^1)} \geq (4i + 1) \cdot (R_{i+1}^{(T_i^0)} + 1)$$

609 Lastly, according to Observation 3.4, there exists a non-zero element in row $i + 1$ at position $i + 1$.
 610 Due to the fact that column $i + 1$ has already been played (according to Item (ii)), we can conclude
 611 that the cumulative utility of row $i + 1$ must be non-zero. Combining this with inductive hypothesis,
 612 it concludes the proof of Item (iv). \square

613 B.2.1 Proof of Theorem 3.13

614 **Proposition B.4.** *There exists a round $T_i^2 > T_i^1$ at which*

- 615 (i) *the strategy profile is $(i + 1, n - i)$ for the first time,*
- 616 (ii) *for all rounds $t \in [T_i^1, T_i^2 - 1]$, the strategy profile is $(i + 1, i + 1)$,*
- 617 (iii) *row $i + 1$ admits cumulative utility $R_{i+1}^{(T_i^2)} \geq (4i + 2) \cdot C_{i+1}^{(T_i^1)}$,*
- 618 (iv) *all rows $k \in [(i + 1) + 1, n - i - 1]$ admit $R_k^{(T_i^2)} = 0$ and all columns $k \in [i + 2, n - (i + 1)]$*
 619 *admit $C_k^{(T_i^2)} = 0$.*

620 *Proof.* We repeat the same reasoning as in the proof of Proposition 3.12. To begin with, we must
 621 establish that the new strategy profile chosen by fictitious play will be $(i + 1, n - i)$.

622 We can infer from both Items (i) and (ii) in Proposition 3.12 that the most recent strategy switch was
 623 made by the row player. Therefore, based on Corollary 3.11, we are certain that the next strategy
 624 switch will be initiated by the column player.

625 In other words, as long as the row player continues to play their current strategy, strategy $(i + 1, i + 1)$
 626 will be played in every subsequent round, which establishes Item (ii). This implies that the column
 627 player's cumulative utility will increase by $e_{i+1}^\top A$. As per Observation 3.3, row $i + 1$ only has
 628 non-zero elements at positions $i + 1$ and $n - i$.

$$e_{i+1}^\top A = [0, \dots, \underbrace{4i + 2}_{i+1}, 0, \dots, 0, \underbrace{4i + 3}_{n-i}, 0, \dots] \quad (6)$$

629 Since strategy $n - i$ has a higher value, a strategy switch in a later time step is certain. Consequently,
 630 there will be a round T_i^2 in which the strategy $(i + 1, n - i)$ will be played for the first time, establishing
 631 Item (i).

632 We must determine the point at which the strategy switch to $(i + 1, n - i)$ will take place. According
 633 to Equation (6), the value of strategy $n - i$ is exactly one greater than the value of strategy $i + 1$.
 634 From the Proposition 3.12, we know that column $n - i$ has a cumulative utility of zero, whereas
 635 column $i + 1$ has a cumulative utility of $C_{i+1}^{(T_i^1)}$. Therefore, it will take precisely $C_{i+1}^{(T_i^1)}$ steps for those
 636 strategies to have equal cumulative utility values. Consequently, the strategy switch will either occur
 637 in that round or the immediate next, as the cumulative utility of column $n - i$ will have surpassed
 638 that of column $i + 1$.

639 In order to proceed with Item (iii), we compute the updated cumulative utility of row $(i + 1)$. As
 640 shown in Equation (6), row $i + 1$ has a value of $(4i + 2)$ at position $i + 1$. Thus, if column $i + 1$
 641 has been played for a minimum of $C_{i+1}^{(T_i^1)}$ rounds, then the cumulative utility of row $i + 1$ satisfies
 642 $R_{i+1}^{(T_i^2)} \geq (4i + 2) \cdot C_{i+1}^{(T_i^1)}$.

643 Lastly, according to Observation 3.4, there exists a non-zero element in column $n - i$ at position
 644 $i + 1$. Due to the fact that row $i + 1$ has already been played (according to Item (ii)), we can conclude
 645 that the cumulative utility of column $n - i$ must be non-zero. Combining this with Proposition 3.12,
 646 it concludes the proof of Item (iv). \square

647 B.2.2 Proof of Theorem 3.14

648 **Proposition B.5.** *There exists a round $T_i^3 > T_i^2$ at which*

- 649 (i) *the strategy profile is $(n - (i + 1), n - i)$ for the first time,*
- 650 (ii) *for all rounds $t \in [T_i^2, T_i^3 - 1]$, the strategy profile is $(i + 1, n - i)$,*
- 651 (iii) *column $n - i$ admits cumulative utility $C_{n-i}^{(T_i^3)} \geq (4i + 3) \cdot R_{i+1}^{(T_i^2)}$,*
- 652 (iv) *all rows $k \in [(i + 1) + 1, n - (i + 1) - 1]$ admit $R_k^{(T_i^3)} = 0$ and all columns $k \in$
 653 $[i + 2, n - (i + 1)]$ admit $C_k^{(T_i^3)} = 0$.*

654 *Proof.* We repeat the same reasoning as in the proof of Proposition 3.12. To begin with, we must
 655 establish that the new strategy profile chosen by fictitious play will be $(n - i - 1, n - i)$.

656 We can infer from both Items (i) and (ii) in Proposition 3.13 that the most recent strategy switch was
 657 made by the column player. Therefore, based on Corollary 3.11, we are certain that the next strategy
 658 switch will be initiated by the row player.

659 In other words, as long as the column player continues to play their current strategy, strategy
 660 $(i + 1, n - i)$ will be played in every subsequent round, which establishes Item (ii). This implies that
 661 the row player's cumulative utility will increase by Ae_{n-i} . As per Observation 3.3, column $n - i$
 662 only has non-zero elements at positions $i + 1$ and $n - (i + 1)$.

$$Ae_{n-i} = [0, \dots, \underbrace{4i + 3}_{i+1}, 0, \dots, 0, \underbrace{4i + 4}_{n-(i+1)}, 0, \dots] \quad (7)$$

Since strategy $n - (i + 1)$ has a higher value, a strategy switch in a later time step is certain. Consequently, there will be a round T_i^3 in which the strategy $(n - i - 1, n - i)$ will be played for the first time, establishing Item (i).

We must determine the point at which the strategy switch to $(n - i - 1, n - i)$ will take place. According to Equation (7), the value of strategy $n - (i + 1)$ is exactly one greater than the value of strategy $i + 1$. From the Proposition 3.13, we know that row $n - (i + 1)$ has a cumulative utility of zero, whereas row $i + 1$ has a cumulative utility of $R_{i+1}^{(T_i^2)}$. Therefore, it will take precisely $R_{i+1}^{(T_i^2)}$ steps for those strategies to have equal cumulative utility values. Consequently, the strategy switch will either occur in that round or the immediate next, as the cumulative utility of row $n - (i + 1)$ will have surpassed that of row $i + 1$.

In order to proceed with Item (iii), we compute the updated cumulative utility of column $n - i$. As shown in Equation (7), column $n - i$ has a value of $(4i + 3)$ at position $i + 1$. Thus, if row $i + 1$ has been played for a minimum of $R_{i+1}^{(T_i^2)}$ rounds, then the cumulative utility of column $n - i$ satisfies $C_{n-i}^{(T_i^3)} \geq (4i + 3) \cdot R_{i+1}^{(T_i^2)}$.

Lastly, according to Observation 3.4, there exists a non-zero element in row $n - (i + 1)$ at position $n - i$. Combining this with the Proposition 3.13 and the fact that column $n - i$ has already been played (according to Item (ii)), we can conclude that the cumulative utility of row $n - (i + 1)$ must be non-zero. This concludes the proof of Item (iv). \square

B.2.3 Proof of Theorem 3.15

Proposition B.6. *There exists a round $T_i^4 > T_i^3$ at which*

- (i) *the strategy profile is $(n - (i + 1), (i + 1) + 1)$ for the first time,*
- (ii) *for all rounds $t \in [T_i^3, T_i^4 - 1]$, the strategy profile is $(n - (i + 1), n - i)$,*
- (iii) *row $n - i - 1$ admits cumulative utility $R_{n-(i+1)}^{(T_i^4)} \geq (4i + 4) \cdot C_{n-i}^{(T_i^3)}$,*
- (iv) *all rows $k \in [(i + 1) + 1, n - (i + 1) - 1]$ admit $R_k^{(T_i^4)} = 0$ and all columns $k \in [(i + 1) + 2, n - (i + 1)]$ admit $C_k^{(T_i^4)} = 0$.*

Proof. We repeat the same reasoning as in the proof of Proposition 3.12. To begin with, we must establish that the new strategy profile chosen by fictitious play will be $(n - (i + 1), (i + 1) + 1)$.

We can infer from both Items (i) and (ii) in Proposition 3.14 that the most recent strategy switch was made by the row player. Therefore, based on Corollary 3.11, we are certain that the next strategy switch will be initiated by the column player.

In other words, as long as the row player continues to play their current strategy, strategy $(n - (i + 1), n - i)$ will be played in every subsequent round, which establishes Item (ii). This implies that the column player's cumulative utility will increase by $e_{n-(i+1)}^\top A$. As per Observation 3.3, row $n - (i + 1)$ only has non-zero elements at positions $(i + 1) + 1$ and $n - (i + 1)$.

$$e_{n-(i+1)}^\top A = [0, \dots, \underbrace{4i + 5}_{(i+1)+1}, 0, \dots, 0, \underbrace{4i + 4}_{n-(i+1)-1}, 0, \dots] \quad (8)$$

Since strategy $(i + 1) + 1$ has a higher value, a strategy switch in a later time step is certain. Consequently, there will be a round T_i^4 in which the strategy $(n - (i + 1), (i + 1) + 1)$ will be played for the first time, establishing Item (i).

We must determine the point at which the strategy switch $(n - (i + 1), (i + 1) + 1)$ will take place. According to Equation (8), the value of strategy $(i + 1) + 1$ is exactly one greater than the value of strategy $n - i$. From the Proposition 3.14, we know that column $(i + 1) + 1$ has a cumulative utility of zero, whereas column $n - i$ has a cumulative utility of $C_{n-i}^{(T_i^3)}$. Therefore, it will take precisely

705 $C_{n-i}^{(T_i^3)}$ steps for those strategies to have equal cumulative utility values. Consequently, the strategy
706 switch will either occur in that round or the immediate next, as the cumulative utility of column
707 $(n - (i + 1), (i + 1) + 1)$ will have surpassed that of column $n - i$.
708 In order to proceed with Item (iii), we compute the updated cumulative utility of row $n - (i + 1)$. As
709 shown in Equation (8), row $n - (i + 1)$ has a value of $(4i + 4)$ at position $n - i$. Thus, if column
710 $n - i$ has been played for a minimum of $C_{n-i}^{(T_i^3)}$ rounds, then the cumulative utility of row $n - (i + 1)$
711 satisfies $R_{n-(i+1)}^{(T_i^4)} \geq (4i + 4) \cdot C_{n-i}^{(T_i^3)}$.
712 Lastly, according to Observation 3.4, there exists a non-zero element in column $(i + 1) + 1$ at position
713 $n - (i + 1)$. Combining this with the Proposition 3.14 and the fact that row $n - (i + 1)$ has already
714 been played (according to Item (ii)), we can conclude that the cumulative utility of column $(i + 1) + 1$
715 must be non-zero. This concludes the proof of Item (iv). \square